Komb: Transparent Combining for Kernel

RS3LAB

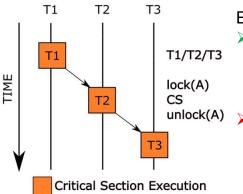
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Existing lock design unable to provide optimal performance

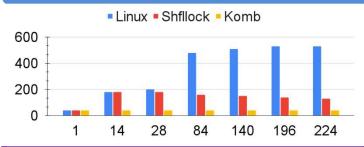
Traditional Design: Move Data to Computation



Existing Locks:

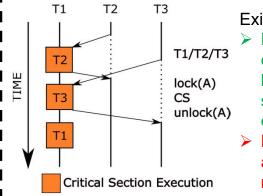
- Reduce contention of lock cache line.
 - Require movement of shared data cache line

Critical Section (CS) Latency



Traditional design unable to reduce CS latency

Alternative Design: Move Computation to Data



Existing Locks:

- Reduce contention of lock and shared data cache line.
- Require application rewrite

Kernel Locks

- Fine-grained locking:
 - > Level locking
- Out-of-order (OOO) unlocks
- Blocking vs non
 - s non- context : ocking ➤ IRQ
 - blocking > If locks
 - Migration

Different

Hard to automatically identify critical section

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Komb Design

Stack Switch

Transfer critical section context using stack.

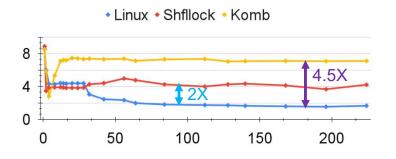
Stack Prefetching

Overlap stack movement with CS execution.

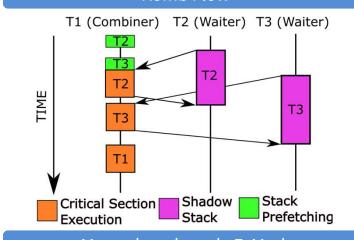
Shadow Stack enables :

- > Async interrupt processing on waiter CPU.
- > Async sleeping of waiter thread.
- ➤ Level locking and OOO unlocking.

Micro-benchmark: Hash-Table



Komb Flow



Macro-benchmark: FxMark



Conclusion: Reduce shared data movement to scale lock performance